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# Physical Network Layer



# Properties of communication networks

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## Constraining factors:

- Transfer rate, (capacity, bandwidth)
- Propagation latency

## Transfer rates:

Morse-telegraph: < 100Bit/sec

Telegraphy: < 150 Bit/sec

Phone: ~ 50Kbit /sec

Serial RS232: ~ 100Kbit/sec

Field bus: few Kbit/sec ... ~ 1Mbit/sec

Ethernet: 10-1000Mbit/sec

High speed networks: >> Gbit/sec

**Latency:**            Satellite connection (2 x 35700 km): ~ 240 ms,  
                          cabel (trans-atlantic) (~ 6.000 km): ~ 20 ms

**Topology:** point-to-point, star, bus, tree, grid, multi-level....



# The Physical Layer Issues

- **Asynchronous serial transmission** (character oriented)
- **Synchronous serial transmission** (bitsynchronization)

- **Bit coding:**
  - NRZ (Non-Return-to-Zero)
  - Manchester Code
  - MFM (Modified-Frequency-Modulation)

- **Modulation and data transmission:**
  - **Base band** Example: Morsetel. / Ethernet
  - **Broad band** Example: Radio, TV, Cabel-TV, Modem
- Modulation: AM, FM

- **Transmissionmedia:**
  - **Fiber** (Multi-Mode, Single-Mode)
  - **Copper** (Twisted Pair, Coaxial)
  - **Radio** (Frequency band)
  - **Satellite** (Geostationary, orbiting)



# How much information can be transferred over a line?

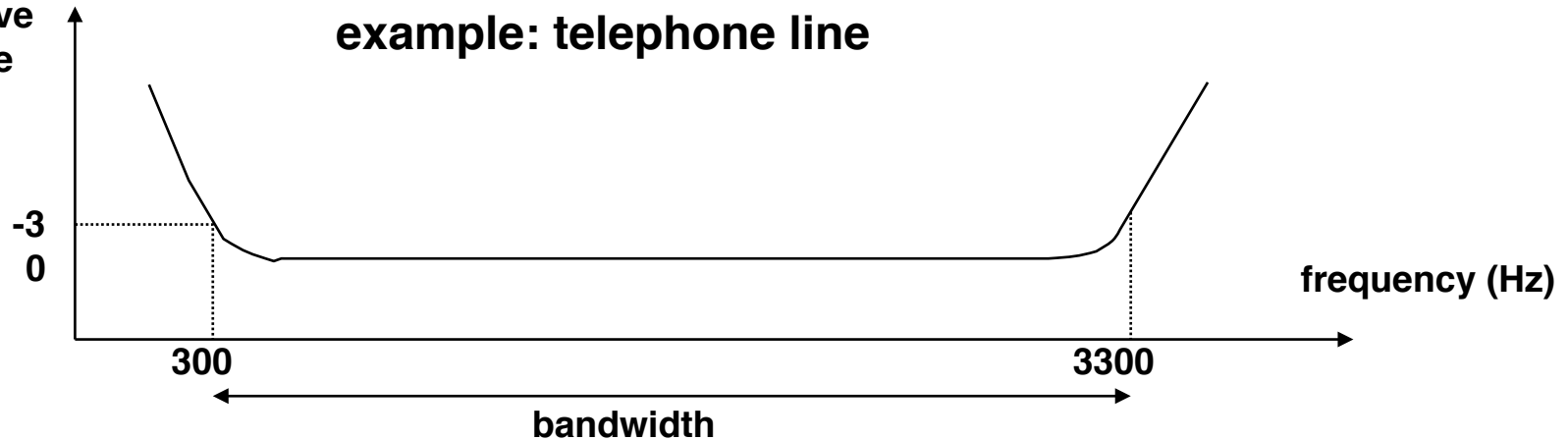
limiting factors:

→ bandwidth of the channel

→ noise

- The bandwidth limits the number of transitions, i.e. the frequency of switching from one signal level to the other
- Noise (informally) limits the ability to distinguish between multiple signal levels

Attenuation of the relative amplitude (in dB)



# Capacity of a channel (Shannon):

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$$C = B \cdot \log_2 (P_s + P_n) / P_n = B \cdot \log_2 (1 + P_s / P_n)$$

**C** : capacity of a channel (measured in Bit/sec (bps))

**P<sub>s</sub>** : signal strength (measured in  $\mu$ W, mW, W)

**P<sub>n</sub>** : noise (measured in  $\mu$ W, mW, W)

**B**: bandwidth

**P<sub>s</sub>/P<sub>n</sub>** : signal-to-noise ratio

**Example:**

**Telephone: bandwidth 3000Hz, signal-to-noise ratio 1000/1**

$$C = 3000 \cdot \log_2 (1+1000) = 3000 \cdot 9,97 = 29900 \text{ Bit/sec (bps)}$$



# Bps and BAUD

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**Bps (Bit/sec) defines a Bit rate (capacity of the channel)**

**BAUD defines the number of symbols. A symbol corresponds to a measurable signal change in the physical transfer medium.**

**Bit/sec is constraint by the channel capacity !**

**BAUD is constraint by the channel bandwidth !**

**Basic methods to increase the bps-Rate at a given BAUD rate of the channel:**

- **distinguish multiple levels**
- **Coding with the smallest number level transitions**



## Coding options (base band)

level  
pulse width  
transitions  
phase

## Bit coding:

NRZ (Non-Return-To-Zero)

Manchester

MFM (Modified-Frequency-Modulation)

## Problems:

synchronization  
number of transitions  
constant/variable frame length

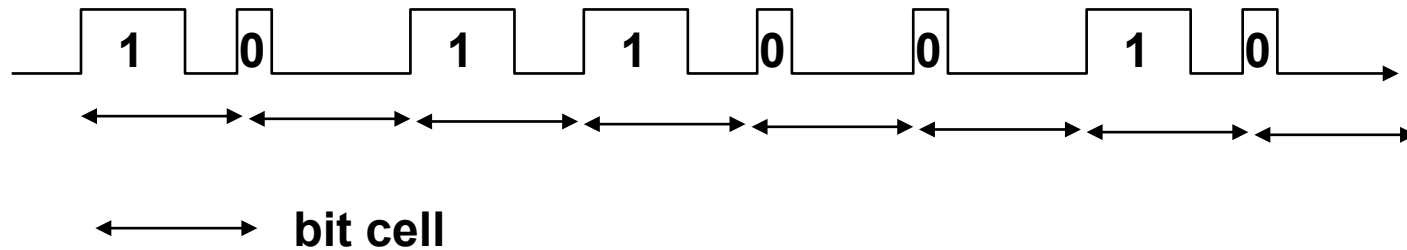


# A simple coding scheme

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**RZ: (Always) Return to Zero (PWM)**

**Example: 1011 0010**



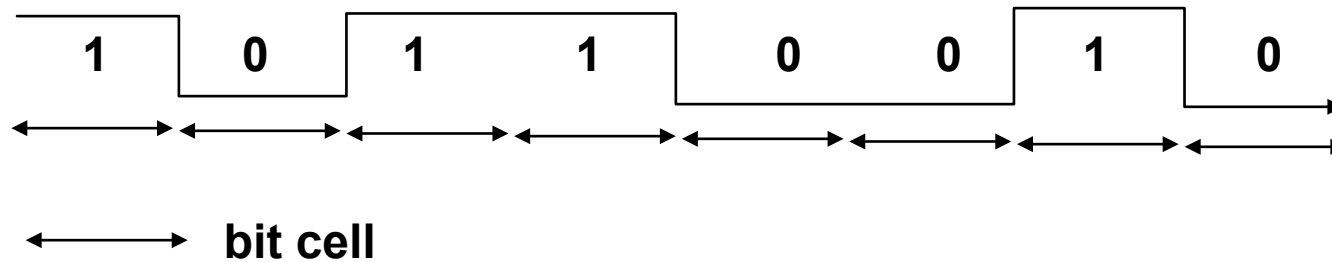


# NRZ Codes

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**NRZ: Non Return to Zero**

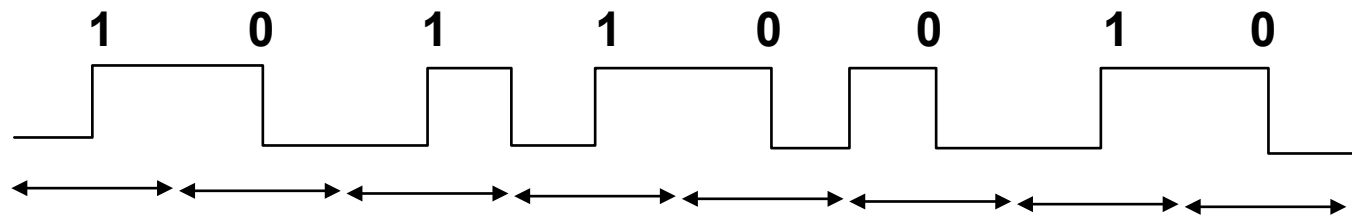
**Example: 1011 0010**



# Manchester Coding

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Example: 1011 0010



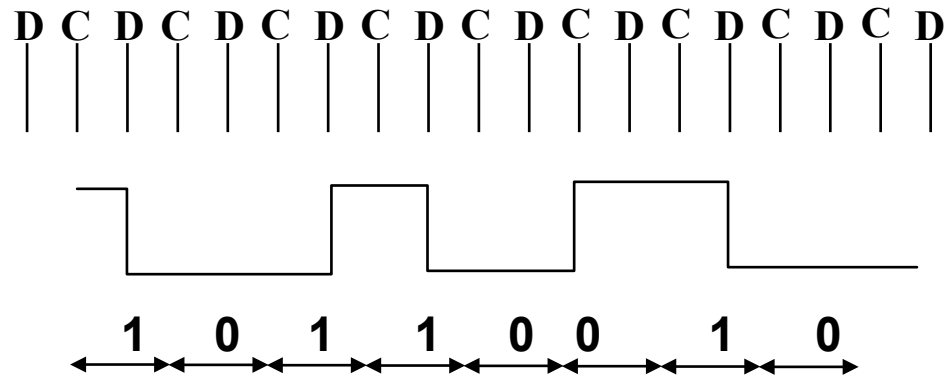
bit cell



# MFM (Modified Frequency Modulation)

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Example: 1011 0010



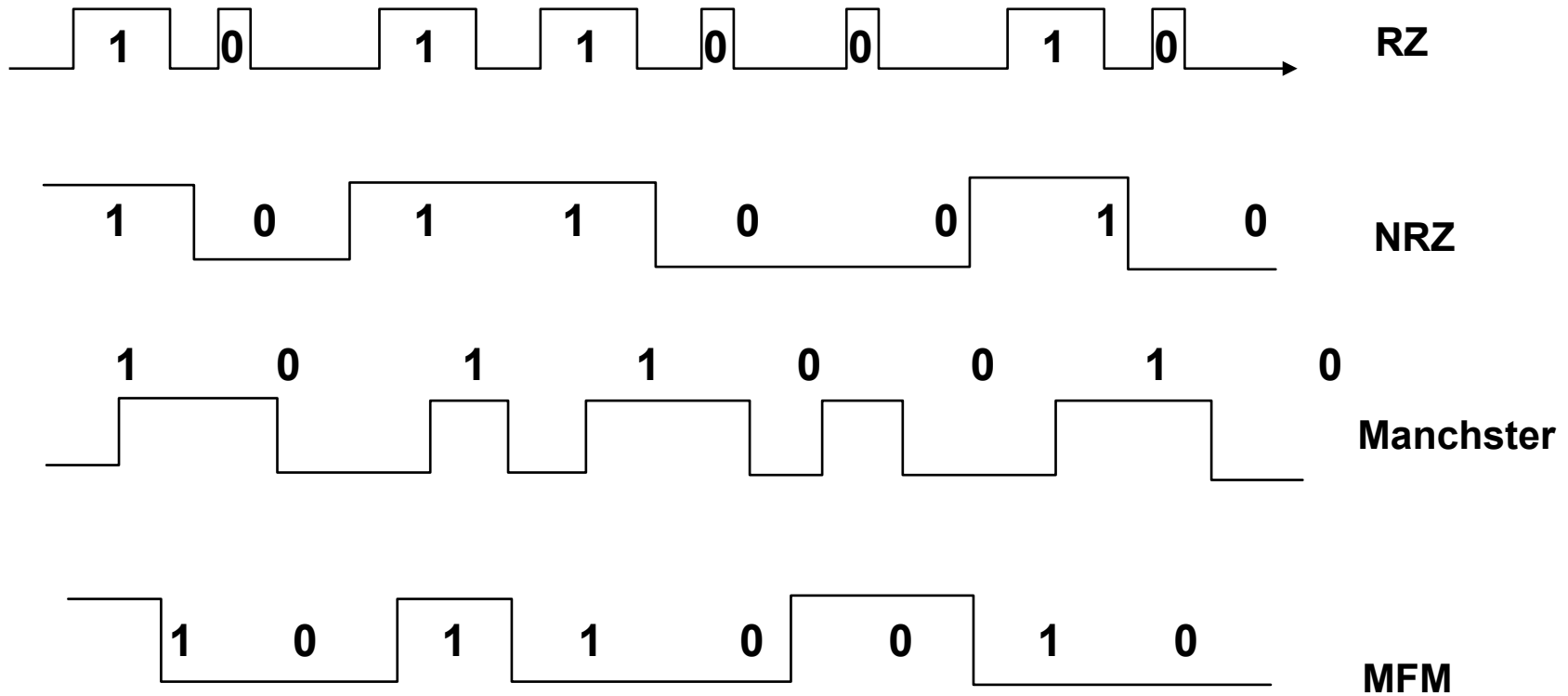
**1: always transition at a data point (D)**

**0: no transition at a data point (D)**

**multiple consecutive "0" : transition at a clock point (C)**



# Comparison of Codes



# Comparison of Codes

Type	Synchronization	transitions/Bit average/max		fixed length
RZ	Y	2	2	Y
NRZ	N	>0,5	1	Y
NRZ*	Y	>0,5	1	N
Manchester	Y	1,5	2	Y
MFM	Y	>0,5	1	Y

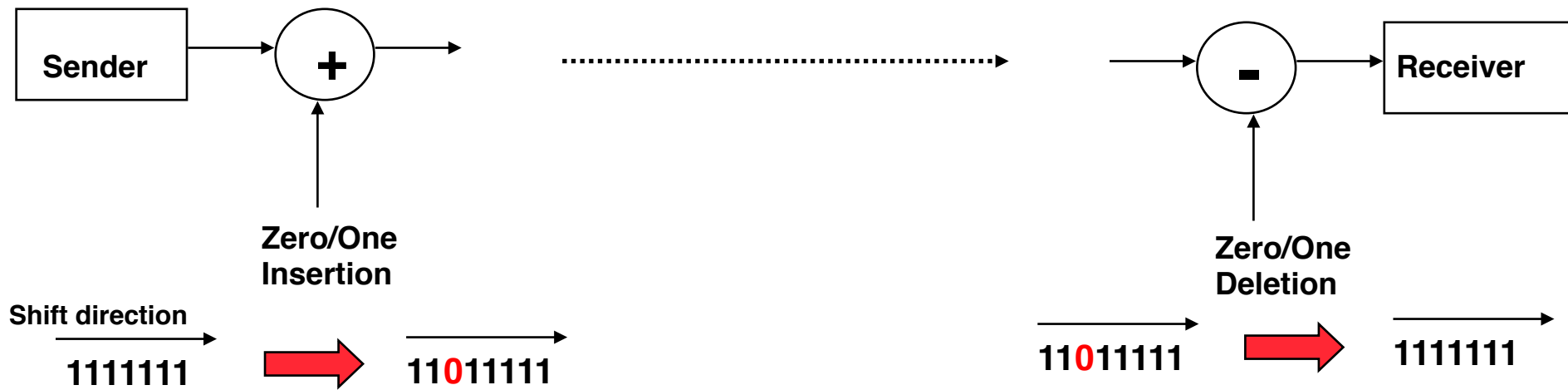
**NRZ\*: NRZ mit Bit Stuffing**



# Bit-Stuffing

When a long sequence of identical values "0" or "1" occurs, bit stuffing inserts a complementary signal level after a fixed specified number of equal signal levels.

Sender transparently inserts stuff bits. The receiver re-establishes the original message by removing the respective stuff bits.



# Bit-Stuffing and framing

## Motivating examples:

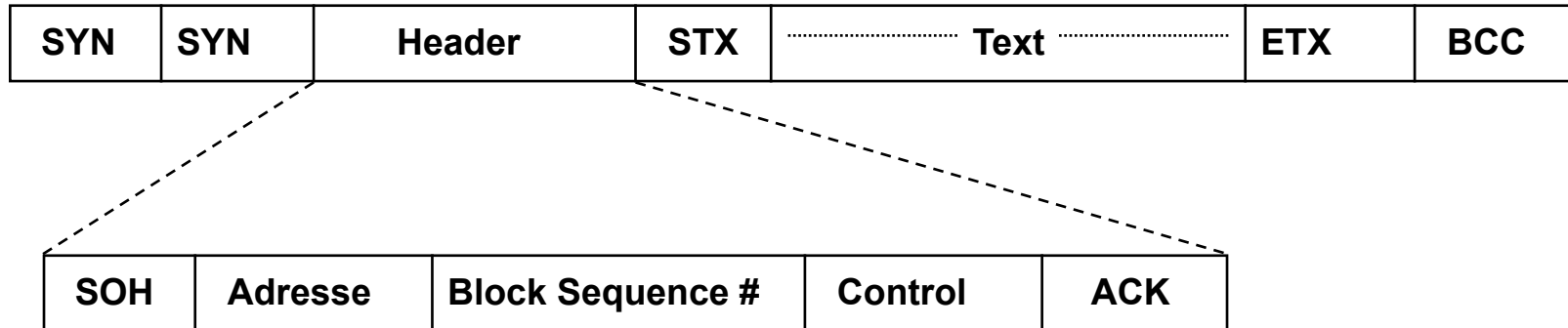
- **character-oriented protocols**  
e.g. IBM BiSync
- **binary protocols**  
e.g. IBM HDLC



# BiSync-Protokoll (IBM)

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Byte-orientiert  
Character-orientiert



## ASCII-Code

<b>SYN :</b>	<b>16<sub>16</sub></b>	
<b>STX :</b>	<b>2<sub>16</sub></b>	<b>Start of Text</b>
<b>ETX:</b>	<b>3</b>	<b>End of Text</b>
<b>BCC:</b>		<b>Error Correcting Code</b>
<b>SOH:</b>	<b>1<sub>16</sub></b>	<b>Start of Header</b>
<b>ACK:</b>	<b>6<sub>16</sub></b>	<b>Acknowledge</b>
<b>NAK:</b>	<b>15<sub>16</sub></b>	<b>Negative Ack.</b>
<b>ENQ:</b>	<b>5<sub>16</sub></b>	<b>Enquire Request.</b>





# Bit stuffing to identify message boundaries

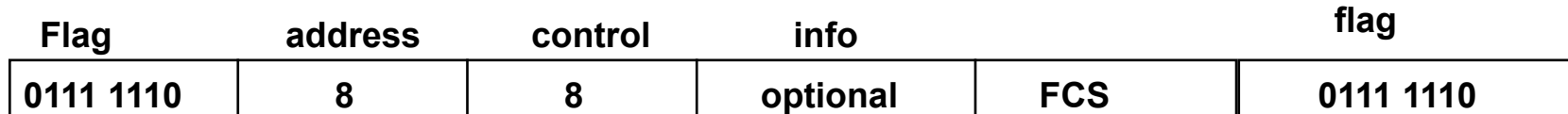
Example: HDLC (High Level Data Link Control)

Problem:

In character-oriented protocols, control and separation characters (STX, ETX, etc.) can be identified easily..  
In bit-oriented protocols any combination of bits as data is possible!

➔ How to identify control information?

HDLC-Frame



I-Frame - Information Frame: data transport  
S-Frame - Supervisor Frame: flow control, e.g. ACK, re-transmission  
U-Frame - Un-numbered Frame: additional control info  
e.g. connect, disconnect

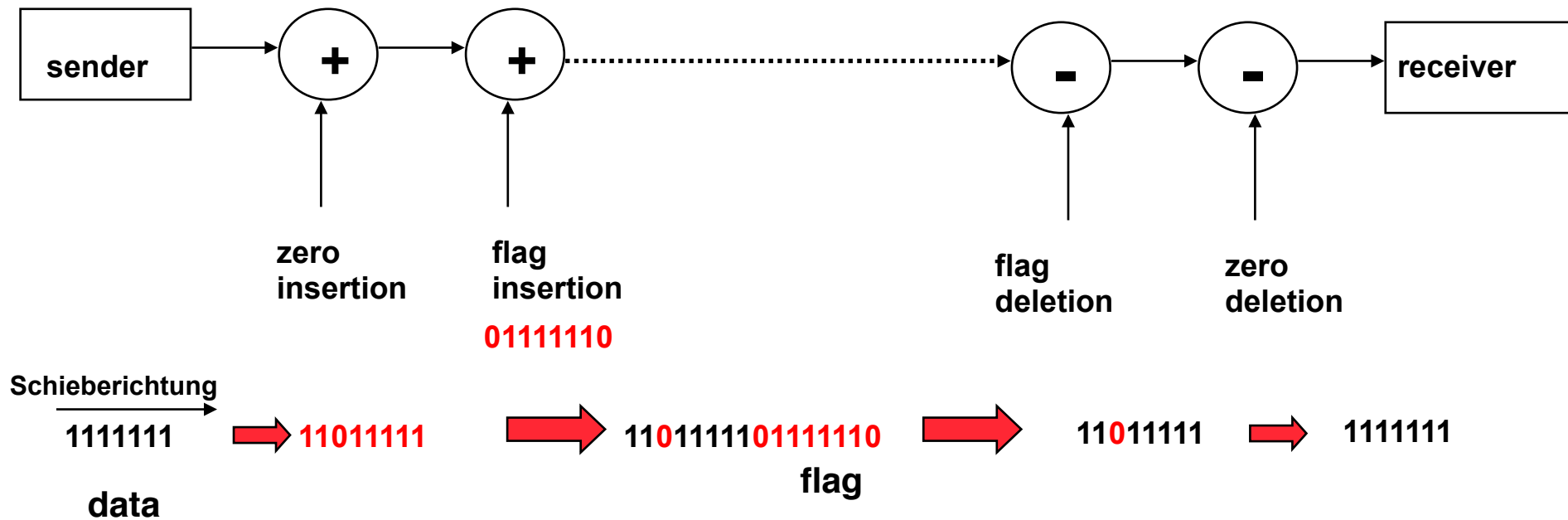
How to distinguish the control info 01111110 from data ?



# Bit stuffing to identify message boundaries

goal: recognizing the flag "01111110".

method: The sender normally inserts a stuff bit "0" after 5 consecutive "1". Therefore there is a max. number of five consecutive "1". The flag is inserted AFTER the bit stuffing stage in the sender and detected and removed before the receiver stuffing stage.



# Bit stuffing

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- **ensures synchronization**
- **needed for framing in binary protocols**
- **exploited for special signals**

